

# The Modern

Cryptography
CookBook

By Anish Nath

# The Modern Cryptography CookBook

Cryptography is for EveryOne. Learn from Crypto Principle to Applied Cryptography With Practical Example

### Anish

This book is for sale at http://leanpub.com/crypto

This version was published on 2018-09-02



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I would like to dedicate this book to my 8-year old son Arjun Nath, Grandfather Sri Rajeshwar Prasad; wife Manisha Prasad; mother Indu Sinha; and all my family members (my father Anil Kumar Sinha; chote papa Sunil Kumar Sinha; choti mummy: Poonam Sinha; and friends). Without them, this would not have been possible.

I would also thanks to Cisco where I work, I got most of the learning from there, and not least the opensource community.

### CONTENTS

# **Contents**

Cryptography	1
Symmetric key Cryptography	4
PUBLIC KEY INFRASTRUCTURES	14
TLSv1.3	28

Cryptography 1

# Cryptography

Cryptography goal is to achieve Confidentiality, Integrity, Identification & Authentication, Non Repudiation

- Data integrity services address the unauthorized or accidental modification of data. The goal is for the receiver of the data to verify that the data has not been altered.
- Confidentiality services restrict access to the content of sensitive data to only those individuals who are authorized to view the data. Confidentiality measures prevent the unauthorized disclosure of information to unauthorized individuals or processes.
- **Identification and authentication** services establish the validity of a transmission, message, and its originator. The goal is for the receiver of the data to determine its origin.
- **Non-repudiation** services prevent an individual from denying that previous actions had been performed. The goal is to ensure that the recipient of the data is assured of the sender's identity.

These four cryptography services is achieved through

- 1. Symmetric key cryptography
- 2. Secure Hash
- 3. Asymmetric (Public-key) cryptography
- 4. Digital Signatures

These **terms** are often used when we discuss cryptography:

- Encryption: The process of converting data from plaintext to ciphertext. Also referred to as enciphering.
- **Decryption**: The process of converting data from ciphertext to plaintext. Also referred to as deciphering.
- **Key**: A parameter that controls the transformation of plaintext into ciphertext or vice versa. Determining the original plaintext data without the key is impossible. Keys can be both public and private. Also referred to as a cryptovariable.
- **Symmetric**: An encryption method whereby a single private key both encrypts and decrypts the data. Also referred to as private or secret key encryption.
- **Asymmetric**: An encryption method whereby a key pair, one private key and one public key, performs encryption and decryption. One key performs the encryption, whereas the other key performs the decryption. Also referred to as public key encryption.

Cryptography 2

• **Digital signature**: A method of providing sender authentication and message integrity. The message acts as an input to a hash function, and the sender's private key encrypts the hash value. The receiver can perform a hash computation on the received message to determine the validity of the message.

- Hash: A one-way function that reduces a message to a hash value. A comparison of the sender's hash value to the receiver's hash value determines message integrity. If the resultant hash values are different, then the message has been altered in some way, provided that both the sender and receiver used the same hash function.
- Digital certificate: An electronic document that identifies the certificate holder.
- Plaintext: A message in its original format. Also referred to as cleartext.
- Ciphertext: An altered form of a message that is unreadable without knowing the key and the encryption system used. Also referred to as a cryptogram.
- Cryptosystem The entire cryptographic process, including the algorithm, key, and key management functions. The security of a cryptosystem is measured by the size of the keyspace and available computational power.
- Cryptanalysis: The science of decrypting ciphertext without prior knowledge of the key or cryptosystem used. The purpose of cryptanalysis is to forge coded signals or messages that will be accepted as authentic signals or messages.
- **Key clustering**: Occurs when different encryption keys generate the same ciphertext from the same plaintext message.
- **Keyspace**: All the possible key values when using a particular algorithm or other security measure. A 40-bit key would have 240 possible values, whereas a 128-bit key would have 2128 possible values.
- Collision: An event that occurs when a hash function produces the same hash value on different messages.
- Algorithm: A mathematical function that encrypts and decrypts data. Also referred to as a cipher.
- Cryptology The science that studies encrypted communication and data.
- **Encoding**: The process of changing data into another form using code.
- Decoding: The process of changing an encoded message back into its original format.
- Transposition: The process of shuffling or reordering the plaintext to hide the original message. Also referred to as permutation. For example, AEEGMSS is a transposed version of MESSAGE.
- **Substitution**: The process of exchanging one byte in a message for another. For example, ABCCDEB is a substituted version of MESSAGE.
- Confusion: The process of changing a key value during each round of encryption. Confusion is often carried out by substitution. Confusion conceals a statistical connection between the plaintext and ciphertext. Claude Shannon first discussed confusion.
- **Diffusion**: The process of changing the location of the plaintext within the ciphertext. Diffusion is often carried out using transposition. Claude Shannon first introduced diffusion.
- Avalanche effect: The condition where any change in the key or plaintext, no matter how minor, will significantly change the ciphertext. Horst Feistel first introduced avalanche effect.

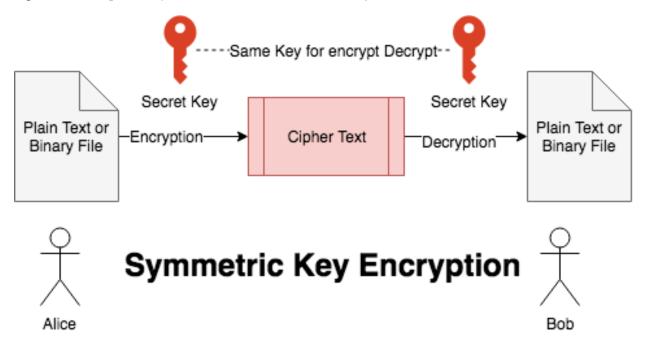
Cryptography 3

• Work factor or work function: The amount of time and resources that would be needed to break the encryption.

- **Trapdoor**: A secret mechanism that allows the implementation of the reverse function in a one-way function.
- One-way function: A mathematical function that can be more easily performed in one direction than in the other.

# Symmetric key Cryptography

**Symmetric key cryptography** is a class of algorithms where **Alice** and **Bob** share a **secret key**. These algorithms are primarily used to achieve **confidentiality**.



#### symmetric Cryptography

Symmetric algorithms are ideally suited for **confidentiality**. To use a symmetric algorithm , Alice transforms a **plaintext** message to **ciphertext** using a symmetric algorithm (AES/Des/3DES..) and a **key**.

Alice transmits the **ciphertext** to **Bob**. Bob uses the **same key** to transform the **ciphertext** back into the **plaintext** 

### Feature of Symmetric key

Symmetric	Integrity	Confidentialitydentification Non Repu-			Key Distri-	
algorithms			& Authen-	diation	bution	
			4! 4!			
			tication			

Symmetric algorithms	Integrity	Confidential	itydentificatio & Authen- tication	n Non Repu- diation	Key Distri- bution
Message Authentica- tion Code	Yes	No	Yes	No	No
key Transport	No	No	No	No	Yes with TTP (Trusted Third Party)

In this chapter we will be discussing below topics related to Symmetric key Cryptography.

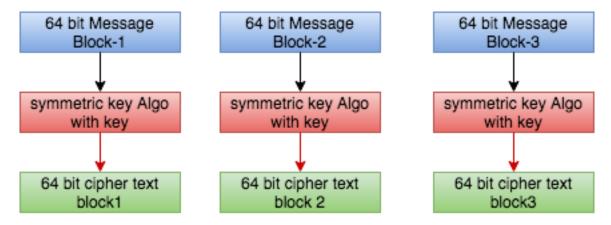
- 1. Modes of Operation
- 2. Authenticated Encrypted Modes
- 3. Initial Vector requirement
- 4. Some Examples Block Cipher Algorithms
  - 1. Data Encryption Standard
  - 2. Triple Data Encryption Standard (TDEA)
  - 3. SkipJack
  - 4. Advanced Encryption Standard
  - 5. RC4
  - 6. BlowFish
- 5. Hash based Symmetric key Algorithms

### **Block ciphers Modes of Operation**

Block ciphers modes provides encryption but not message integrity

### ECB:

The Electronic Codebook (ECB), The message is divided into blocks, and each block is encrypted separately, this is the simplest encryption and decryption mode.



# **ECB Mode**

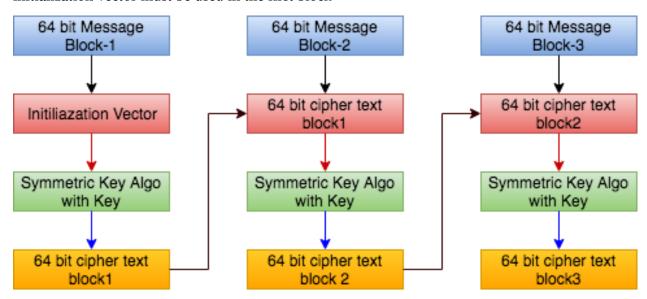
ecb

### **Example Ciphers**

AES-128-ECB, AES-192-ECB, AES-256-ECB, BF-ECB, CAMELLIA-128-ECB, CAMELLIA-192-ECB, CAMELLIA-256-ECB, CAST5-ECB, DES-ECB, IDEA-ECB, RC2-ECB, RC5-ECB, SEED-ECB, AES-128-ECB, AES-192-ECB, AES-256-ECB, BF-ECB, CAMELLIA-128-ECB, CAMELLIA-192-ECB, CAMELLIA-256-ECB, CAST5-ECB, DES-ECB, IDEA-ECB, RC2-ECB, RC5-ECB, SEED-ECB

#### **CBC**

Cipher Block Chaining (CBC) mode an IV-based encryption scheme , each block of plaintext is XORed with the previous ciphertext block before being encrypted. This way, each ciphertext block depends on all plaintext blocks processed up to that point. To make each message unique, an initialization vector must be used in the first block



### **CBC Mode**

ecb

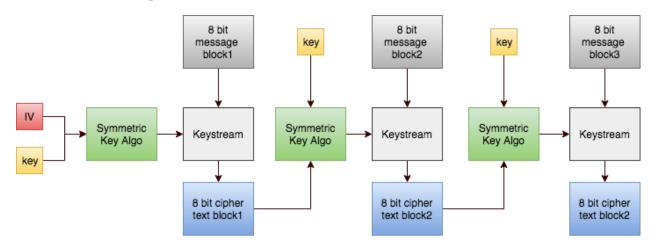
#### **Example Ciphers**

AES-128-CBC, AES-128-CBC-HMAC-SHA1, AES-128-CBC-HMAC-SHA256, AES-192-CBC, AES-256-CBC, AES-256-CBC-HMAC-SHA1, AES-256-CBC-HMAC-SHA256, AES128 => AES-128-CBC, AES192 => AES-192-CBC, AES256

=> AES-256-CBC, BF => BF-CBC, BF-CBC, CAMELLIA-128-CBC, CAMELLIA-192-CBC, CAMELLIA-256-CBC, CAMELLIA-256-CBC, CAMELLIA-256-CBC, CAMELLIA-128-CBC, CAMELLIA-128-CBC, CAMELLIA-256-CBC, CAST
=> CAMELLIA-128-CBC, CAMELLIA192 => CAMELLIA-192-CBC, CAMELLIA-256-CBC, CAST
=> CAST5-CBC, CAST-cbc => CAST5-CBC, CAST5-CBC, DES => DES-CBC, DES-EDE-CBC, DES-EDE-CBC, DES-EDE3-CBC, DESS-EDE-CBC, DES-EDE3-CBC, DESS-EDE3-CBC, DESS-EDE3-CBC, RC2-CBC, RC2-CBC, RC3-CBC, RC3-CBC,

### **CFB**

Cipher Feedback (CFB) mode is a confidentiality mode that features the feedback of successive ciphertext segments into the input blocks of the forward cipher to generate output blocks that are exclusive-ORed with the plaintext to produce the ciphertext, and vice versa. The CFB mode requires an IV as the initial input block



### **CFB Mode**

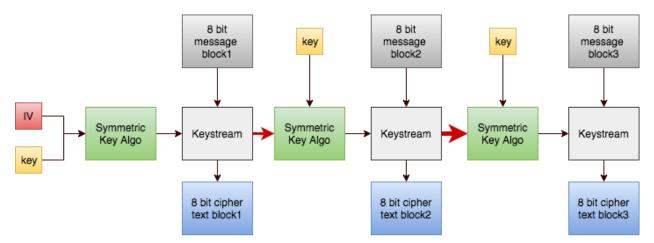
cfb

#### **Example Ciphers**

AES-128-CFB, AES-128-CFB1, AES-128-CFB8, AES-192-CFB, AES-192-CFB1, AES-192-CFB8, AES-256-CFB, AES-256-CFB, AES-256-CFB1, AES-256-CFB3, BF-CFB, CAMELLIA-128-CFB, CAMELLIA-128-CFB1, CAMELLIA-128-CFB8, CAMELLIA-192-CFB3, CAMELLIA-192-CFB3, CAMELLIA-256-CFB3, CAMELLIA-256-CFB1, CAMELLIA-256-CFB1, CAMELLIA-256-CFB3, CAMELLIA-256-CFB3, DES-CFB3, DES

#### **OFB**

The Output Feedback (OFB) mode makes a block cipher into a synchronous stream cipher. It generates keystream blocks, which are then XORed with the plaintext blocks to get the ciphertext



### **OFB Mode**

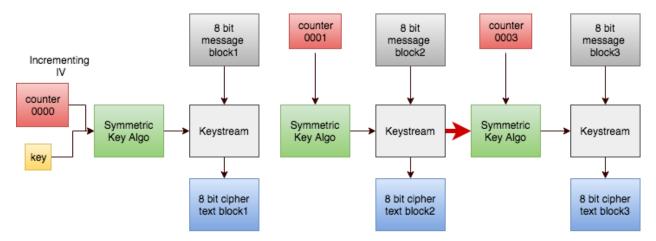
ofb

### **Example Ciphers**

AES-128-OFB, AES-192-OFB, AES-256-OFB, BF-OFB, CAMELLIA-128-OFB, CAMELLIA-192-OFB, CAMELLIA-256-OFB, CAST5-OFB, DES-EDE-OFB, DES-EDE3-OFB, DES-OFB, IDEA-OFB, RC2-OFB, RC5-OFB, SEED-OFB, AES-128-OFB, AES-192-OFB, AES-256-OFB, BF-OFB, CAMELLIA-128-OFB, CAMELLIA-192-OFB, CAMELLIA-256-OFB, CAST5-OFB, DES-EDE-OFB, DES-EDE3-OFB, DES-OFB, IDEA-OFB, RC2-OFB, RC5-OFB, SEED-OFB

### **CTR**

The CTR mode has similar characteristics to OFB, but also allows a random access property during decryption. CTR mode is well suited to operate on a multi-processor machine where blocks can be encrypted in parallel. Furthermore, it does not suffer from the short-cycle problem that can affect OFB.



### **CTR Mode**

ctr

### **Example Ciphers**

AES-128-CTR, AES-192-CTR, AES-256-CTR, AES-128-CTR, AES-192-CTR, AES-256-CTR

### **XTS**

An IV-based encryption scheme, the mode works by applying a tweakable blockcipher (secure as a strong-PRP) to each n-bit chunk. For messages with lengths not divisible by n, the last two blocks are treated specially. The only allowed use of the mode is for encrypting data on a block-structured storage device. The narrow width of the underlying PRP and the poor treatment of fractional final blocks are problems. More efficient but less desirable than a (wide-block) PRP-secure blockcipher would be.

### **Example Ciphers**

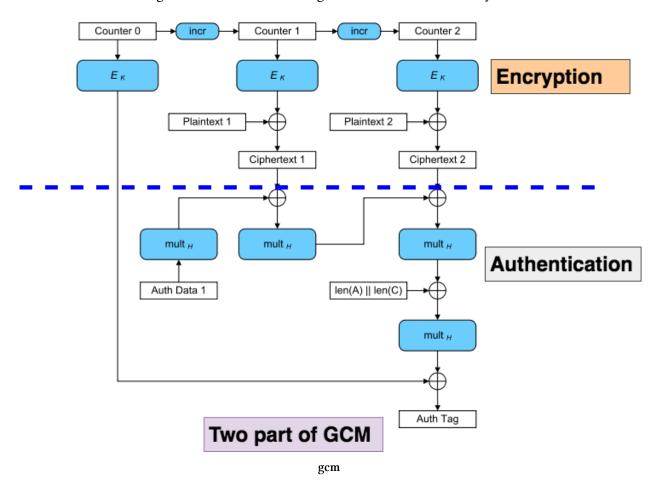
AES-128-XTS, AES-256-XTS, AES-128-XTS, AES-256-XTS

### **Authenticated encryption Modes**

Authenticated encryption provides both encryption and message integrity

### **GCM**

Galois/Counter Mode (GCM) A nonce-based AEAD scheme that combines CTR mode encryption and a GF(2128)-based universal hash function. Good efficiency characteristics for some implementation environments. Good provably-secure results assuming minimal tag truncation. Attacks and poor provable-security bounds in the presence of substantial tag truncation. Can be used as a nonce-based MAC, which is then called GMAC. Questionable choice to allow nonces other than 96-bits. Recommend restricting nonces to 96-bits and tags to at least 96 bits. Widely standardized and used.



### **Example Ciphers**

id-aes128-GCM,id-aes192-GCM,id-aes256-GCM,id-aes128-GCM,id-aes192-GCM,id-aes256-GCM

### CCM

CCM mode (Counter with CBC-MAC), A nonce-based AEAD scheme that combines CTR mode encryption and the raw CBC-MAC. Inherently serial, limiting speed in some contexts. Provably secure, with good bounds, assuming the underlying blockcipher is a good PRP. Ungainly construction that demonstrably does the job. Simpler to implement than GCM. Can be used as a nonce-based MAC. Widely standardized and used.

### **Example Ciphers**

id-aes128-CCM, id-aes192-CCM, id-aes256-CCM, id-aes128-CCM, id-aes192-CCM, id-aes256-CCM

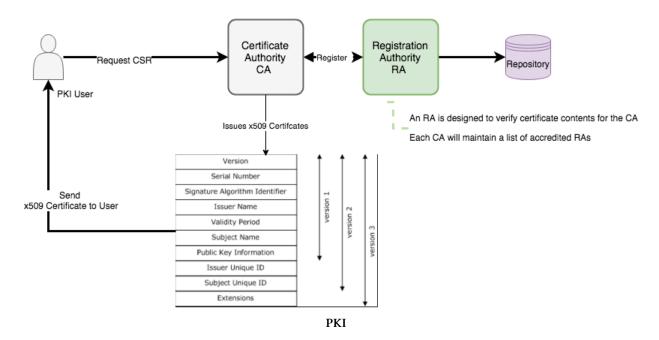
### **Initialization Vectors**

Generation of Initial Vectors - The CBC, CFB, and OFB modes require an initialization vector as input, in addition to the plaintext. An IV must be generated for each execution of the encryption operation, and the same IV is necessary for the corresponding execution of the decryption operation - For the CBC and CFB modes, the IVs must be unpredictable

### **PUBLIC KEY INFRASTRUCTURES**

A public key infrastructure (PKI) binds public keys to entities, enables other entities to verify public key bindings, and provides the services needed for ongoing management of keys in a distributed system

### **PKI COMPONENTS**



### **Certification Authority (CA)**

The CA issues a public key certificate for each identity, confirming that the identity has the appropriate credentials.

#### Main Functions of CA

- Verify the CSR request
- Issue the Certificate (Create and SIgn them)
- Attach CRL for Certificate revocation
- Publish it's current (Expired Certificate ) and CRL's

PUBLIC KEY INFRASTRUCTURES 15

### **Registration Authority (RA)**

An RA is designed to verify certificate contents for the CA. Each CA will maintain a list of accredited RAs; that is a list of RAs determined to be trustworthy

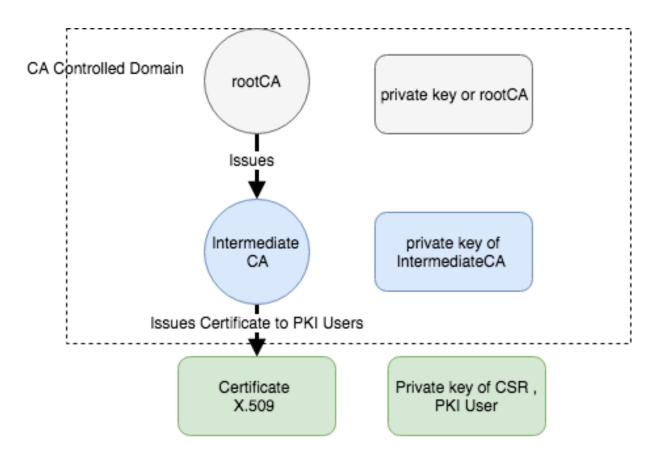
### **PKI Users**

PKI Users are organizations or individuals that use the PKI, but do not issue certificates.

**Main Functions of PKI user** - Generate the Certificate Signing Request. - Maintain the Certificate obtain by the CA.

### **PKI Architecture**

CAs may be linked in a number of ways. Most enterprises that deploy a PKI will choose either a **mesh** or a **hierarchical** architecture, This is an example of **hierarchical** Structure



# **Hierarchical PKI Arcitecture**

PKI

### **Generating hierarchical CA structure**

• rootCA will generate self signed certificate and key with longer validity

```
openssl genrsa -des3 -out rootCA.key 4096
openssl req -new -x509 -days 3650 -key rootCA.key -out rootCA.crt
```

• intCA will generate CSR and get it signed with rootCA and set Validity for Longer Year.

```
openssl genrsa -des3 -out intCA.key 4096 openssl req -new -key intCA.key -out intCA.csr
```

intCA Submitted the CSR information to rootCA to get it signed

PUBLIC KEY INFRASTRUCTURES 17

```
Enter pass phrase for intCA.key:
You are about to be asked to enter information that will be incorporated
into your certificate request.
What you are about to enter is what is called a Distinguished Name or a DN.
There are quite a few fields but you can leave some blank
For some fields there will be a default value,
If you enter '.', the field will be left blank.
Country Name (2 letter code) [AU]:
State or Province Name (full name) [Some-State]:KA
Locality Name (eg, city) []:IN
Organization Name (eg, company) [Internet Widgits Pty Ltd]: IntermediateCA Issuers
Organizational Unit Name (eg, section) []:intermediateCA Issue of myOrg
Common Name (e.g. server FQDN or YOUR name) []:intCA
Email Address []:
Please enter the following 'extra' attributes
to be sent with your certificate request
A challenge password []:
An optional company name []:
```

• Sign **intermediate**CA with rootCA private key, control the serial Number to **specific** pattern (Usability)

```
openssl x509 -req -days 365 -in intCA.csr -CA rootCA.crt -CAkey rootCA.key -set_se\
rial 1111111 -out intCA.crt
Signature ok
subject=/C=AU/ST=KA/L=IN/O=IntermediateCA Issuers/OU=intermediateCA Issue of myOrg\
/CN=intCA
Getting CA Private Key
Enter pass phrase for rootCA.key:
```

• PKI (End)Users will create the CSR and send to CA to get it signed

```
openssl genrsa -des3 -out client.key 2048 openssl req -new -key client.key -out client.csr
```

End user submit the CSR information to get it signed by the CA, the CA used intermediate CA to sign the CSR

```
Enter pass phrase for client.key:
You are about to be asked to enter information that will be incorporated
into your certificate request.
What you are about to enter is what is called a Distinguished Name or a DN.
There are quite a few fields but you can leave some blank
For some fields there will be a default value,
If you enter '.', the field will be left blank.
Country Name (2 letter code) [AU]:
State or Province Name (full name) [Some-State]: My Org
Locality Name (eg, city) []:IN
Organization Name (eg, company) [Internet Widgits Pty Ltd]:8gwifi
Organizational Unit Name (eg, section) []:Crypto
Common Name (e.g. server FQDN or YOUR name) []:8gwifi.org
Email Address []:
Please enter the following 'extra' attributes
to be sent with your certificate request
A challenge password []:
An optional company name []:
```

• Sign with Intermediate CA,set the exipry date to 1 or 2 year Max, and generate a serial number for this

```
openssl x509 -req -in client.csr -days 530 -CA intCA.crt -CAkey intCA.key -CAcreat\ eserial -out client.crt
```

#### The CSR getting signed

```
Signature ok subject=/C=AU/ST=My Org/L=IN/O=8gwifi/OU=Crypto/CN=8gwifi.org Getting CA Private Key Enter pass phrase for intCA.key:
```

• View the Client Certificate Information , The issuer is IntermediateCA

#### Verify Certificate

Verify Client certificate with Full CA chain

```
openssl verify -verbose -CAfile <(cat intCA.crt rootCA.crt) client.crt client.crt: OK
```

Verify intermediateCA belongs to rootCA chain

```
openssl verify -verbose -CAfile rootCA.crt intCA.crt intCA.crt: OK
```

### **PKI Data Structure**

Two Basic data structures used in PKI - X.509 Public Key Certificates - CRL Certificate Revocation List - Attribute Certificates

### X.509

There are ten common fields, six mandatory and four optional.

The mandatory fields are: serial number, signature algorithm identifier, certificate issuer name, certificate validity period, public key Info, and the subject name

X.509 Attributes	Description
Version	v2 or v3
Serial Number	CA Assigned Serial Number to the
Signature	Certificate Indicate which digital Signature
	Algorithms ex: SHA-256 with RSA
	Encryption
Issuer	Contains x.500 DN
Validity	Certificate Expiry Dates
Subject	Contains the DN of the Holder of private
	Key Corresponding to the public key in the
	certificate

X.509 Attributes	Description
Subject public key Information	Optional Parameter, Algorithm Identifiers
Issuer UniqueID and Subject UniqueID	ID of the Issues only in v2 and v3
Extensions	Optional only in v3 Certificate version
subjectType	Indicate whether a subject is CA or ENtity
Names and identity information	c=US; o=8gwifi; ou=Crypto; cn=8gwifi.org
Key Attributes	Specifies relevant attributes whether it
	used for key transport, or be used to verify
	a digital signature
Policy Information	Policies related to Certificate
Certificate Extensions	extension identifier, a criticality flag, and
	extension value

### Reader Note what is self Signed Certificate

In self signed certificate the issuer and subject are same :)

```
Issuer: C=AU, ST=KA, L=IN, O=rootCA Issuers, OU=rootCA, CN=rootCA Subject: C=AU, ST=KA, L=IN, O=rootCA Issuers, OU=rootCA, CN=rootCA
```

### **CRL**

The CRL contains the following fields:

CRL Fields	Description
Version	Optional Version, Default 2
Signature	algorithm identifier for the digital signature algorithm used
	by the CRL issuer to sign the CRL
Issuer	X.500 DN of the CRL issuer
This update	Issue date of the CRL
Next update	Next CRL Issue date
Revoked Certificates	Structured List of Revoked Certificates
CRL Extensions	Additional Information About the CRL

#### Pem format of CRL

```
----BEGIN X509 CRL----
```

### Implement CERTIFICATION REVOCATION LIST

• Make a directory for a CRL:

```
mkdir -p /etc/pki/crl
```

• Create an **index file**, **the CRL Database** with the following command:

```
touch /etc/pki/crl/index.txt
```

• Create a file for the CRL number. This file should contain the text 00 only.

```
echo 00 > /etc/pki/crl/crl_number
```

• crl\_openssl.conf: create and write the following contents into a crl\_openssl.conf file.

```
cat <<EOF > crl_openssl.conf
# OpenSSL configuration for CRL generation
[ ca ]
default ca = CA default
                                 # The default ca section
[ CA default ]
database = /etc/pki/crl/index.txt
crlnumber = /etc/pki/crl_number
                                    # how long to certify for
default_days
             = 365
                                # how long before next CRL
default_crl_days= 30
default_md
            = default
                               # use public key default MD
                                # keep passed DN ordering
preserve
           = no
[ crl_ext ]
# CRL extensions.
# Only issuerAltName and authorityKeyIdentifier make any sense in a CRL.
# issuerAltName=issuer:copy
authorityKeyIdentifier=keyid:always,issuer:always
EOF
```

• Generate CRL file

```
openssl ca -gencrl -keyfile intCA.key -cert intCA.crt -out intCA.crl -config crl_\ openssl.conf
```

• View CRL file

Added with required CRL field, at this time no certificate is Invoked

#### • Revoke the Certificate

```
openssl ca -revoke client.crt -keyfile intCA.key -cert intCA.crt -config crl_opens\
sl.conf
Using configuration from crl_openssl.conf
Enter pass phrase for intCA.key:
Adding Entry with serial number AC12C39820C69327 to DB for /C=AU/ST=Some-State/O=I\
nternet Widgits Pty Ltd/CN=8gwifi.org
Revoking Certificate AC12C39820C69327.
Data Base Updated
```

### • Again revoking the same client certificate will through an error Already Revoked

```
openssl ca -revoke client.crt -keyfile intCA.key -cert intCA.crt -config crl_opens\
sl.conf
Using configuration from crl_openssl.conf
Enter pass phrase for intCA.key:
ERROR:Already revoked, serial number AC12C39820C69327
```

#### Regenerate the CRL list

23

```
openssl ca -gencrl -keyfile intCA.key -cert intCA.crt -out intCA.crl -config crl_o\penssl.conf
Using configuration from crl_openssl.conf
Enter pass phrase for intCA.key:
```

· View CRL file

#### One revoked certificate is Added in the CRL entry

### Adding CRL distribution point

• Edit the file crl\_openssl.conf and point out the PEM and DER

```
crlDistributionPoints=@crl_section
[crl_section]
URI.1 = https://8gwifi.org/intCA.crl
URI.2 = https://8gwifi.org/intCA.der
```

• CA role, once CA sign the file it will use -extfile crl\_openssl.conf to locate the crlDistributionPoints

```
openssl x509 -req -in client.csr -days 530 -CA intCA.crt -CAkey intCA.key -CAcreat\ eserial -out client.crt -extfile crl_openssl.conf
```

 Once certificate is Signed, the x.509 certificate will have X509v3 extensions that contains CRL Distribution Points

### **OCSP**

OCSP stands for the Online Certificate Status Protocol and is one way to validate a certificate status. It is an alternative to the CRL, certificate revocation list.

### **OCSP Request** tbsRequest Version requesterName requestList hashAlgorithm issuerNameHash CA issuerKeyHash serialNumber Client OCSP Responder OCSP singleRequestExtension requestExtensions optional Signature signatureAlgorithm signature certs **OCSP**

#### The OCSP process is very simple:

- 1. Client receives the certificate
- 2. Client sends OCSP request to the OCSP server and it query by the serial number of the certificate
- 3. OCSP response with a certificate status Good, Revoked or Unknown

### Working Demo

• Get the certificate you want to verify 8gwifi.org1

```
openssl s_client -servername 8gwifi.org -connect 8gwifi.org:443 2\times1 < /dev/null |\ sed -n '/----BEGIN/,/----END/p' > 8gwifi.pem
```

• Build the certificate chain

```
openssl s_client -servername 8gwifi.org -connect 8gwifi.com:443 -showcerts 2 \% 4 \dev/null > cacert.pem
```

edit the file cacert.pem and add necessary chain certificate

• Determine the ocsp URI

```
openssl x509 -noout -ocsp_uri -in 8gwifi.pem
http://ocsp.int-x3.letsencrypt.org
```

• Invoke the openssl ocsp client

```
openssl ocsp -no_nonce -issuer cacert.pem -cert 8gwifi.pem -VAfile cacert.pem -te\
xt -url http://ocsp.int-x3.letsencrypt.org/ -header Host ocsp.int-x3.letsencrypt.org\
-respout ocsptest
```

The OCSP Response

¹https://8gwifi.org

```
OCSP Request Data:
    Version: 1 (0x0)
    Requestor List:
        Certificate ID:
          Hash Algorithm: sha1
          Issuer Name Hash: 7EE66AE7729AB3FCF8A220646C16A12D6071085D
          Issuer Key Hash: A84A6A63047DDDBAE6D139B7A64565EFF3A8ECA1
          Serial Number: 03FF3497BFA5D45C36C511809F9FD5F28C20
OCSP Response Data:
    OCSP Response Status: successful (0x0)
    Response Type: Basic OCSP Response
    Version: 1 (0x0)
    Responder Id: C = US, O = Let's Encrypt, CN = Let's Encrypt Authority X3
    Produced At: Aug 6 08:59:00 2018 GMT
    Responses:
    Certificate ID:
      Hash Algorithm: sha1
      Issuer Name Hash: 7EE66AE7729AB3FCF8A220646C16A12D6071085D
      Issuer Key Hash: A84A6A63047DDDBAE6D139B7A64565EFF3A8ECA1
      Serial Number: 03FF3497BFA5D45C36C511809F9FD5F28C20
    Cert Status: good
    This Update: Aug 6 08:00:00 2018 GMT
    Next Update: Aug 13 08:00:00 2018 GMT
    Signature Algorithm: sha256WithRSAEncryption
         39:3d:96:78:44:9f:03:29:bc:83:35:32:1a:d7:6d:05:f9:59:
         63:9e:52:6f:06:8e:9f:74:d1:f9:aa:18:2b:e2:13:61:5a:d1:
         ad:7a:67:9e:2b:a1:12:83:92:92:f3:c0:dc:4b:2a:ee:96:85:
         f0:5b:39:30:2f:17:ed:20:a3:ae:de:c1:41:e1:26:8d:70:c5:
         fe:79:9c:37:7b:b7:75:93:61:f7:5f:8b:7f:6f:99:7a:5a:19:
         a4:e7:4b:41:ad:e5:92:71:44:11:75:67:68:0d:0c:b6:be:ef:
         70:a0:a7:c6:fa:6e:06:08:5a:7c:2e:f0:41:7a:55:a3:21:74:
         89:2c:e5:f9:ab:58:5c:97:1d:89:a8:65:a3:be:f7:0a:e5:5c:
         4f:a9:61:f2:04:d5:f2:18:6b:74:e7:b5:c8:12:db:9c:70:89:
         e7:c6:e3:43:70:18:41:d6:4b:a9:15:94:13:4b:00:75:d2:2a:
         fe:fb:e4:a8:cf:e5:aa:56:d6:e1:91:55:06:d1:33:43:d9:4b:
         82:a6:bc:10:a1:42:d0:e2:49:fe:18:08:44:d4:a7:4f:b6:3f:
         00:95:72:11:d7:e3:14:eb:6b:51:7a:e7:c1:40:42:2e:da:c4:
         be:1a:ce:8c:48:f2:03:ed:c0:93:19:c9:26:93:1e:f0:d2:56:
         bc:70:39:db
Response verify OK
8gwifi.pem: good
        This Update: Aug 6 08:00:00 2018 GMT
```

Next Update: Aug 13 08:00:00 2018 GMT

### **TLSv1.3**

TLS stands for Transport Layer Security<sup>2</sup> and is the successor to SSL (Secure Sockets Layer). TLS allows client/server applications to communicate over the Internet in a way that is designed to prevent eavesdropping, tampering, and message forgery. This section mostly focus on TLSv 1.3. TLS 1.3 was defined in RFC<sup>3</sup> 8446<sup>4</sup> in August 2018. It is based on the earlier TLS 1.2 specification.

### **TLS Histrory and Development**

Protocol	Published	
SSL 1.0	Unpublished	
SSL 2.0	1995	
SSL 3.0	1996	
TLS 1.0	1999	
TLS 1.1	2006	
TLS 1.2	2008	
TLS 1.3	2018	

TLS1.3 has been over eight years since the last encryption protocol update, but the final version of TLS 1.3 has now been published as of August 2018 Image Ref<sup>5</sup>



<sup>&</sup>lt;sup>2</sup>https://en.wikipedia.org/wiki/Transport\_Layer\_Security

<sup>&</sup>lt;sup>3</sup>https://en.wikipedia.org/wiki/Request\_for\_Comments\_(identifier)

<sup>4</sup>https://tools.ietf.org/html/rfc8446

<sup>&</sup>lt;sup>5</sup>https://datatracker.ietf.org/doc/rfc8446/

### **Major Differences from TLS 1.2**

ssl						Expression +
No.	Time	Source	Destination	Protocol	Length Info	
	4 0.000509	127.0.0.1	127.0.0.1	TLSv1.3	299 Client Hello	
	6 0.008516	127.0.0.1	127.0.0.1	TLSv1.3	2299 Server Hello, Change Cipher Spec, Application Data, Application Data,	Application Data, App
	8 0.009185	127.0.0.1	127.0.0.1	TLSv1.3	146 Change Cipher Spec, Application Data	
	9 0.009564	127.0.0.1	127.0.0.1	TLSv1.3	321 Application Data	
	10 0.009759	127.0.0.1	127.0.0.1	TLSv1.3	321 Application Data TLSv1.3	
	12 7.048517	127.0.0.1	127.0.0.1	TLSv1.3	90 Application Data	
ssi						Expression +
No.	Time	Source	Destination	Protocol	Length Info	
	4 0.000508	127.0.0.1	127.0.0.1	TLSv1.2	377 Client Hello	
	6 0.007982	127.0.0.1	127.0.0.1	TLSv1.2	2142 Server Hello, Certificate, Server Key Exchange, Server Hello Done	
	8 0.008630	127.0.0.1	127.0.0.1	TLSv1.2	159 Client Key Exchange, Change Cipher Spec, Encrypted Handshake Message	TLSv1.2
	9 0.008886	127.0.0.1	127.0.0.1	TLSv1.2	292 New Session Ticket, Change Cipher Spec, Encrypted Handshake Message	12341.2

From the Wireshark packet capture, its clearly visible the TLSv.1.3, the number of packets is being reduced this offer better speed in TLS v1.3, and some of the major changes from TLS1.2 as follows

• The list of supported symmetric encryption algorithms has been pruned of all algorithms that are considered legacy. Those that remain are all **Authenticated Encryption with Associated Data** (AEAD) algorithms.

```
    Handshake Protocol: Client Hello
    Handshake Type: Client Hello (1)
    Length: 224
    Version: TLS 1.2 (0x0303)
    Random: 3ab2121f850c067cc82af4c012f5ac3ddb1e7e6369c1eb91...
    Session ID Length: 32
    Session ID: 19785f578273db9bfa2848077cc3bb2698fc6260eec96e51...
    Cipher Suites Length: 8
    ▼ Cipher Suites (4 suites)
        Cipher Suite: TLS_AES_256_GCM_SHA384 (0x1302)
        Cipher Suite: TLS_CHACHA20_POLY1305_SHA256 (0x1303)
        Cipher Suite: TLS_AES_128_GCM_SHA256 (0x1301)
        Cipher Suite: TLS_EMPTY_RENEGOTIATION_INFO_SCSV (0x00ff)
```

- The cipher suite concept has been changed to separate the authentication and key exchange mechanisms from the record protection algorithm (including secret key length) and a hash to be used with both the key derivation function and handshake message authentication code (MAC).
- A zero round-trip time (0-RTT) mode was added, saving a round trip at connection setup for some application data, at the cost of certain security properties. IMP 0-rtt should be avoided, there are proven replay attack has been found
- All handshake messages after the **ServerHello** are now **encrypted**. The newly introduced EncryptedExtensions message allows various extensions previously sent in the clear in the ServerHello to also enjoy confidentiality protection.

```
icr paytoau (2233 bytes)
▼ Secure Sockets Layer
  ▼ TLSv1.3 Record Layer: Handshake Protocol: Server Hello
       Content Type: Handshake (22)
       Version: TLS 1.2 (0x0303)
       Length: 122

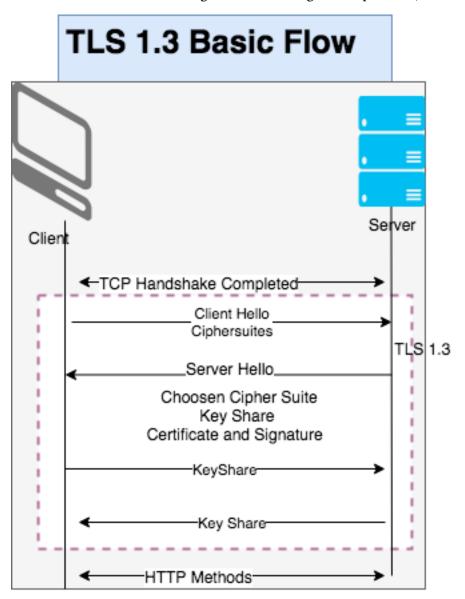
▼ Handshake Protocol: Server Hello

         Handshake Type: Server Hello (2)
         Length: 118
         Version: TLS 1.2 (0x0303)
         Random: e1b018a3eaecf9dd1496daeab292841ae8c65cf811c70b6d...
         Session ID Length: 32
         Session ID: 1913c49cf29a36d42a71dc5347227f31e2d9894215ed63af...
         Cipher Suite: TLS_AES_256_GCM_SHA384 (0x1302)
         Compression Method: null (0)
         Extensions Length: 46
       ▶ Extension: supported_versions (len=2)
       ▶ Extension: key_share (len=36)
  ▶ TLSv1.3 Record Layer: Change Cipher Spec Protocol: Change Cipher Spec
    TLSv1.3 Record Layer: Application Data Protocol: http-over-tls
       Opaque Type: Application Data (23)
                                                                        Encrypted
       Version: TLS 1.2 (0x0303)
       Length: 23
       Encrypted Application Data: 48fc81eb1428148de99ef7da96b1e1ee1e34326c9ae2e5
  ▶ TLSv1.3 Record Layer: Application Data Protocol: http-over-tls
    TLSv1.3 Record Layer: Application Data Protocol: http-over-tls
    TLSv1.3 Record Layer: Application Data Protocol: http-over-tls
```

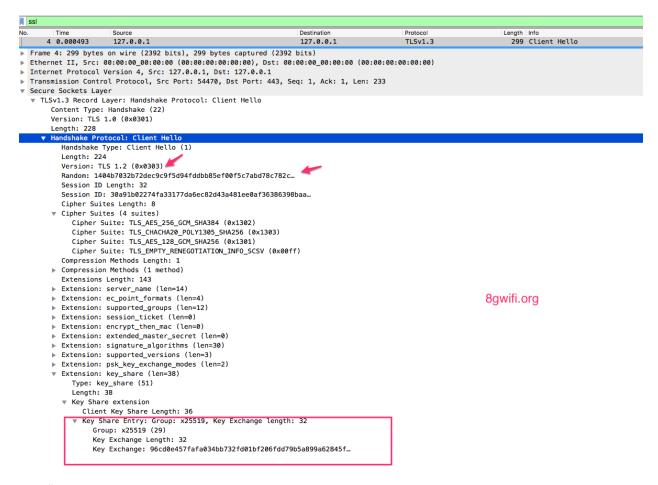
- Static **RSA** and **Diffie-Hellman** cipher suites have been **removed**; all public-key based key exchange mechanisms now provide forward secrecy.
- The key derivation functions have been redesigned.
- The handshake state machine has been significantly restructured to be more consistent and to remove superfluous messages such as ChangeCipherSpec (except when needed for middlebox compatibility).
- Elliptic curve algorithms are now in the base spec, and new signature algorithms, such as EdDSA, are included. TLS 1.3 removed point format negotiation in favor of a single point format for each curve.
- The TLS 1.2 version **negotiation mechanism has been deprecated** in favor of a version list in an extension. This increases compatibility with existing servers that incorrectly implemented version negotiation.
- Session resumption with and without server-side state and the PSK-based ciphersuites of earlier versions of TLS have been replaced by a single new PSK exchange

### TLS 1.3 Handshake

The handshake can be thought of as having three phases (indicated in the diagram below)



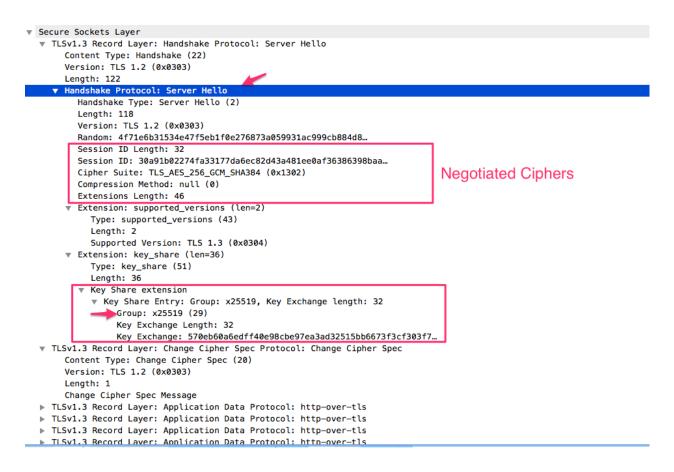
- 1. In the first phase, the client sends the ClientHello message, which contains
  - random nonce (ClientHello.random);
  - protocol versions;
  - symmetric cipher/HKDF hash pairs; either a set of Diffie-Hellman key shares (in the "key\_share" extension
  - A set of pre- shared key labels (in the "pre\_shared\_key" extension or both;
  - And potentially additional extensions.



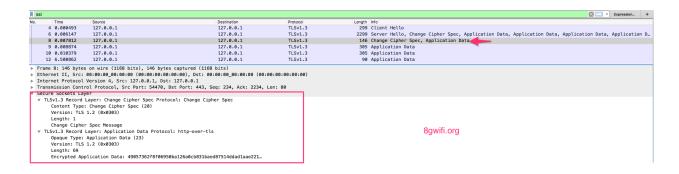
The "key\_share" extension contains the endpoint's cryptographic parameters.In TLSv1.3 the client selects a "group� that it will use for key exchange.

The PSK: If clients offer "pre\_shared\_key" without a "psk\_key\_exchange\_modes" extension, servers abort the handshake and used to negotiate the identity of the pre-shared key to be used with a given handshake in association with PSK key establishment

1. The server processes the ClientHello and determines the appropriate cryptographic parameters for the connection. It then responds with its own ServerHello which indicates the negotiated connection parameters. The combination of the ClientHello and the ServerHello determines the shared keys



1. Upon receiving the server's messages, the client responds with its Authentication messages, namely Certificate and CertificateVerify (if requested), and Finished.



### **Ciphersuites**

OpenSSL has implemented support for five TLSv1.3 ciphersuites as follows:

1. TLS\_AES\_256\_GCM\_SHA384

```
2. TLS_CHACHA20_POLY1305_SHA256
```

- 3. TLS\_AES\_128\_GCM\_SHA256
- 4. TLS\_AES\_128\_CCM\_8\_SHA256
- 5. TLS\_AES\_128\_CCM\_SHA256

```
openssl ciphers -v | grep TLSv1.3

TLS_AES_256_GCM_SHA384 TLSv1.3 Kx=any Au=any Enc=AESGCM(256) Mac=AEAD

TLS_CHACHA20_POLY1305_SHA256 TLSv1.3 Kx=any Au=any Enc=CHACHA20/POLY1305(256) Mac=\
AEAD

TLS_AES_128_GCM_SHA256 **TLSv1.3** Kx=any Au=any Enc=AESGCM(128) Mac=AEAD
```

### Start the TLS1.3 server in openssl

The forthcoming openssl 1.1.1-pre9 (beta) release has included support for TLSv1.3.

```
openssl version
OpenSSL 1.1.1-pre9 (beta) 21 Aug 2018

openssl command to start the tls1.3 server

openssl s_server -accept 443 -tls1_3 -ciphersuites TLS_AES_256_GCM_SHA384 -key key.\
pem -cert cert.pem
Using default temp DH parameters
ACCEPT
```

### **Connect to TLS1.3**

The openssl command to connect to tlsv1.3.

```
openssl s_client -connect 127.0.0.1:443
----
New, TLSv1.3, Cipher is TLS_AES_256_GCM_SHA384
Server public key is 4096 bit
Secure Renegotiation IS NOT supported
Compression: NONE
Expansion: NONE
No ALPN negotiated
Early data was not sent
```

Verify return code: 18 (self signed certificate)

---

Post-Handshake New Session Ticket arrived:

SSL-Session:

Protocol : TLSv1.3

Cipher: TLS\_AES\_256\_GCM\_SHA384

Session-ID: 2BC1AB6B0BE58B527AE4CAEFEABC6D9654094BC1F4D529E5F3F0912A80C97001

Session-ID-ctx:

Resumption PSK: EA4A8E23B397F4F822B770C0922F47F7A66F6A7AA2F2DC4B94B961941AA87ACD\

611AC293259EFB130887F9A2D02AC89E

PSK identity: None

PSK identity hint: None

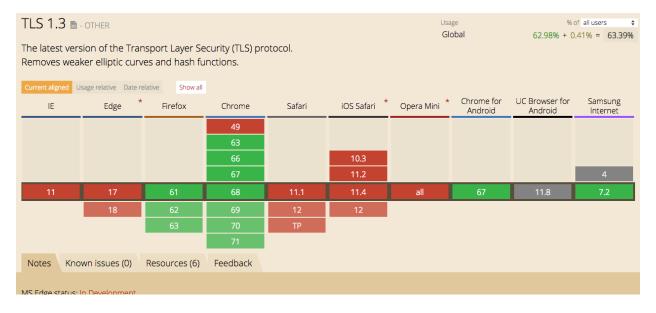
SRP username: None

TLS session ticket lifetime hint: 7200 (seconds)

TLS session ticket:

### **Browser Support**

Checkout the browser compatibility for TLS 1.3 here: https://caniuse.com/#feat=tls1-3



# **Further Reading**

RFC  $^6$  8446  $^7$  Cloudflare TLS 1.3 and Q&A  $^8$ 

 $<sup>{}^{6}</sup>https://en.wikipedia.org/wiki/Request\_for\_Comments\_(identifier) \\ {}^{7}https://tools.ietf.org/html/rfc8446 \\ {}^{8}https://blog.cloudflare.com/tls-1-3-overview-and-q-and-a/}$