

Subject: MCR-10053, Fix exercise_disk and rdisk_ Event Wait Channel Processing

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Problem Description

Multics Change Ticket: <http://multics-trac.swenson.org/ticket/163> describes an issue with the exercise_disk command and the rdisk_ i/o module. The issue is that the code that does i/o using ioi_\$connect does not always call ipc_\$block. Not calling ipc_\$block results in an event message remaining in the ECT for the event channel, which may or may not get processed for the next ioi_\$connect call. Unprocessed event messages may accumulate, overflowing the process directory in which the area that holds the ECT lives.

Each time ioi_\$connect is called, an interrupt from the IOM will be sent to Multics indicating the status of the i/o. The status can indicate an error or success, but this interrupt is converted into a wakeup event message for the event wait channel associated with the IOI device. This event wait channel is established in the initial rcp_\$attach (or rcp_priv_\$attach) call made to attach the device for i/o. The expectation is that each event message will be read from the ECT during a call to ipc_\$block and then freed from the area in which the ECT lives.

The ECT is the event channel table, and it lives in an per-ring area in the process directory. A field in each ring's stack header points to the ECT area. When the first event channel is created in a ring, ipc_real_ calls an internal procedure, find_ectp, which checks the value of the ect_ptr in the stack header for the ring. If it is null, it calls ipc_util_\$create_ect to create the ECT. This latter entry allocates an area for the ECT (ect_area) in the system free area. It then calls define_area_ to convert the original area into an extensible area. An extensible area can grow beyond a segment. The area manager knows how to add new segments to the area (which is called a multisegment or extensible area) when there is not enough room in the area for a new allocation. An internal procedure of ipc_util_\$create_ect allocates an ect_header in the extensible area, initializes the ECT, and sets a pointer to this ect_header in the stack header (for the current ring).

Ring-0 IPC allocates ITT (Interprocess Transmission Table) messages in this area when the user-ring IPC calls hcs_\$read_events. These ITT messages are threaded into either an event call or event wait message thread in the ECT by user-ring IPC. When ipc_\$block is called in the user ring, the code retrieves the first event message associated with the event channel passed in a structure parameter to ipc_\$block. If no call to ipc_\$block is made (nor the event channel drained with ipc_\$drain_chn) then the event message stays in the ECT.

The completion status of IOI i/o is reflected in a status queue in the IOI workspace and updated automatic by the interrupt termination support in ioi_masked\$interrupt. This status queue is directly accessible to user-ring code, and can therefore be checked after a call to ioi_\$connect. It is possible that this status can get updated after a call to ioi_\$connect and before a call to ipc_\$block to (possibly) wait for the termination status. The code in exercise_disk and rdisk_ has an optimization — it checks the status for completion and errors and if already updated, skips the call to ipc_\$block. Unfortunately, this has the side effect of leaving the IOI event in the ECT where it will get picked up the next time ipc_\$block is called.

If this happens, the event message returned from `ipc_$block` will NOT correspond to the immediately preceding call to `ioi_$connect`, but to an earlier call. And the other side effect is that the event messages may accumulate in the ECT.

The emulator for Multics has the ability to support synchronous i/o (as well as, in a newer release, asynchronous i/o). With synchronous i/o, the emulator posts the terminate interrupt before returning from the “connect” signal sent by Multics to begin the i/o. This results in completion status being available very rapidly – possibly just after the call to `ioi_$connect` and before the call to `ipc_$block`. So sometimes, completion status is available when `exercise_disk` or `rdisk_` checks the status word. And in this case, the code doesn’t call `ipc_$block`, with the consequence that event messages accumulate in the ECT.

`exercise_disk` takes a long time to run—even for smaller 451 disks. In a recent run it took 13 hours to run (I was running a debug version of the emulator, which is slower than a non-debug version, so this time might have been longer as a result). Well before `exercise_disk` finishes, when synchronous emulator i/o is used, the number of event messages that accumulate in the ECT causes the quota on the process directory to overflow (my process directory quota is 2000 records). `exercise_disk` does a lot of i/o operations and therefore one can very quickly overflow quota in the process directory. Each new ITT message allocated in the ECT area can potentially grow the area. And because the area is an extensible area, it adds new segments as needed. These segments consume more and more quota in the process directory.

If asynchronous i/o is enabled in the emulator, no quota overflow occurs. While it is possible, due to the speed of the underlying hardware and the emulator that some completion status is posted before the call to `ipc_$block` is made, and thus some event messages remain in the ECT when they shouldn’t, the number must be far less or perhaps even zero, because quota does not appear to grow.

The problematic code in `exercise_disk` is:

```
do while (again);                                /* I/O loop */

    completion.st = "0";                          /* initialize status entry */
    completion.run = "1";

    call ioi_$connect (devx, dcw_offset, code);    /* Start I/O */
    if code ^= 0 then call io_err ("0"); /* didn't get away from the starting line */

    do while (^completion.st & completion.run);    /* while connected and no status */

        call ipc_$block (addr (wait_list), addr (event_info), code); /* wait for completion */
        if code ^= 0 then call io_err ("0");      /* No loiterers?? */

    end;
```

As can be seen, the call to `ipc_$block` is skipped if there completion status or if the i/o is no longer in progress. This is the case where the event message is not extracted from the ECT and therefore accumulates.

Proposed Changes

The proposed change is to add a flag, called `block`, which is set to false at each `ioi_$connect` iteration, and which is set to true when `ipc_$block` is called.

The above code is changed to:

```
do while (again);                                /* I/O loop */

    completion.st = "0"b;                          /* initialize status entry */
    completion.run = "1"b;

    call ioi_$connect (devx, dcw_offset, code); /* Start I/O */
    if code ^= 0 then call io_err ("0"b);        /* didn't get away from the starting line */

    called_block = "0"b;
    do while (^completion.st & completion.run); /* while connected and no status */

        call ipc_$block (addr (wait_list), addr (event_info), code); /* wait for completion */
        if code ^= 0 then call io_err ("0"b); /* No loiterers?? */
        called_block = "1"b;

    end;
    if ^called_block then
        call ipc_$drain_chn (event_info.chan_id, (0));
```

The event wait channel used by `exercise_disk.pl1` and `rdisk_.pl1` is the same event channel used by RCP when attaching the device. That event channel becomes the IOI event channel used for i/o. While there is no evidence that the same race condition described above that can leave event messages in the ECT for this event channel occurs during RCP attachment, in the event that any event messages are left in the channel after the RCP attachment, a defensive additional call to `ipc_$drain_chn` is added at the tail end of the “mount” procedure — which performs the RCP attachment. This guarantees that prior to the first `ioi_$connect` call to perform i/o, there are no event messages in the ECT for the event channel.

The compare_ascii output for the change follows:

```
Inserted in B:
B109          dcl      called_block          bit (1);
Preceding:
A104          dcl      code                  fixed bin (35);

Inserted in B:
B1369         dcl      ipc_$drain_chn        entry (fixed bin(71), fixed bin(35));
Preceding:
A1363         dcl      ioi_$set_status       entry (fixed bin, fixed bin (18), fixed bin (8), fixed bin (35));

Inserted in B:
B1978         called_block = "0"b;
Preceding:
A1971         do while (^completion.st & completion.run); /* while connected and no status */

A1975
A1976         end;
Changed by B to:
B1983         called_block = "1"b;
B1984
B1985         end;
B1986         if ^called_block then          /* drain events in case we didn't call ipc_$block
*/
B1987         call ipc_$drain_chn (wait_list.ev_chan, (0));
```

```

Inserted in B:
B2507          /* drain event wait channel of any events that might have been read during rcp attachment */
B2508          call ipc_$drain_chn (ev_chan, (0));
B2509
Preceding:
A2496          detach: proc;

Comparison finished: 6 differences, 42 lines.
r 14:58 1.027 13

```

A similar, corresponding change will be made to `rdisk_`, which has the same problem. The `compare_ascii` output for the change follows:

```

Inserted in B:
B385          dcl  ipc_$drain_chn                      entry (fixed bin(71), fixed bin(35));
Preceding:
A380          dcl  iox_$err_no_operation                entry options (variable);

Inserted in B:
B1625          dcl  called_block                        bit (1);                      /* keep track of whether we
picked up terminate event */
B1626
Preceding:
A1619          true_len = min (4 * block_len, data_left);          /* set true amount to xmit
*/

Inserted in B:
B1757          called_block = "0"b;
Preceding:
A1749          do while (^completion.st & completion.run); /* while connected and no
status */

A1754
A1755          end;
A1756
Changed by B to:
B1763          called_block = "1"b;                      /* remember that we've
retrieved termination event */
B1764          end;
B1765          if ^called_block then                      /* drain events in case we
didn't call ipc_$block */
B1766          call ipc_$drain_chn (wait_list.ev_chan, (0));
B1767

Inserted in B:
B2794          /* drain event wait channel of any events that might have been read during
rcp attachment */
B2795          call ipc_$drain_chn (ev_chan, (0));
B2796
Preceding:
A2783          detach:

Comparison finished: 8 differences, 40 lines.
r 15:00 1.752 43

```

Documentation

The proposed changes do not affect the function and descriptions of either `exercise_disk` or `rdisk_` and therefore no documentation changes are necessary.

Testing

The changes to `exercise_disk` have already been tested by running the program on a 451 disk. Prior to the changes, running `exercise_disk` provoked a `record_quota_overflow` signal. After the changes, the program completed successfully.

`rdisk_` will be tested by performing a `reload_volume`, which uses `rdisk_` to restore a volume from volume backup tapes.

Version History

Date	Revision	Author	Comment
2019-04-21	1.0	Eric Swenson	Initial revision of the MCR.
2019-04-21	1.1	Eric Swenson	Added description of ECT area explain how the ECT can grow beyond one segment, and contribute to a <code>record_quota_overflow</code> .
2019-05-18	1.2	Eric Swenson	Updated <code>compare_ascii</code> output for <code>exercise_disk.pl1</code> based on review comments from Gary. Added <code>rdisk_.pl1</code> <code>compare_ascii</code> output. Added description and justification for additional call to <code>ipc_\$drain_chn</code> called at tail end of mount procedure.