The Linux Kernel Scheduler

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Topics

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- Scheduling classes and policies
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CPU Scheduler

- Shares CPU between tasks.
- Selects next task to run (on each CPU) when:
  - Running task terminates
  - Running task sleeps (waiting for an event)
  - A new task created or sleeping task wakes up
  - Running task’s timeslice is over
- Goals:
  - Be *fair*
  - Timeslice based on task priority
  - Low task response time
  - High (task completion) throughput
  - Balance load among CPUs
  - Power efficient
  - Low overhead of running scheduler’s code
- Work with mainframes, servers, desktops/laptops, embedded/phones.
The O(1) scheduler

- 140 priorities. 0-99: RT tasks and 100-139: User tasks.
- Each CPU’s runqueue had 2 arrays: Active and Expired.
- Each array had 140 entries, one per priority level.
- Each entry was a linked list serviced in FIFO order.
- Bitmap (of 140 bits) used to find which priority lists aren’t empty.
- Timeslices were assigned to tasks based on these priorities.
- Expired tasks moved from Active to Expired array.
- Once Active array was empty, the arrays were swapped.
- Enqueue and dequeue of tasks and next task selection done in constant time.
- Replaced by CFS in Linux 2.6.23 (2007).
The O(1) scheduler (Cont..)

CPU Active array
- Priority 0
- Priority 1
- ...
- Priority 99
- ...
- Priority 139

CPU Expired array
- Priority 0
- Priority 1
- ...
- Priority 99
- ...
- Priority 139

Real Time task priorities

User task priorities
Current scheduler design

- Introduced by Ingo Molnar in 2.6.23 (2007).
- Scheduling policies within scheduling classes.
- Scheduling class with higher priority served first.
- Task can migrate between CPUs, scheduling policies and scheduling classes.
Scheduling Classes

- Represented by following structure:

```c
struct sched_class {
    const struct sched_class *next;
    void (*enqueue_task) (struct rq *rq, struct task_struct *p, int flags);
    void (*dequeue_task) (struct rq *rq, struct task_struct *p, int flags);
    struct task_struct * (*pick_next_task) (struct rq *rq, struct task_struct *prev, struct rq_flags *rf);
    ...
};
```

![Diagram showing the scheduling classes network](image-url)
Schedule()

- Picks the next runnable task to run on a CPU.
- Searches for a task starting from the highest priority class (stop).
- Helper: for_each_class().
- pick_next_task():

```c
again:
    for_each_class(class) {
        p = class->pick_next_task(rq, prev, rf);
        if (p) {
            if (unlikely(p == RETRY_TASK))
                goto again;
            return p;
        }
    }

/* The idle class should always have a runnable task: */
BUG();
```
Scheduling classes and policies

- **Stop**
  - No policy
- **Deadline**
  - SCHED_DEADLINE
- **Real Time**
  - SCHED_FIFO
  - SCHED_RR
- **Fair**
  - SCHED_NORMAL
  - SCHED_BATCH
  - SCHED_IDLE
- **Idle**
  - No policy
Sched class: STOP

- Highest priority class.
- Only available for SMP (stop_machine() isn’t used in UP).
- Can preempt everything and is preempted by nothing.
- Mechanism to stop running everything else and run a specific function on CPU(s).
- No scheduling policies.
- One kernel thread per CPU: “migration/N”.
- Used by task migration, CPU hotplug, RCU, ftrace, clockevents, etc.
Sched class: Deadline (DL)

- Highest priority tasks in the system.
- Scheduling policy: SCHED_DEADLINE.
- Implemented with red-black tree (self balancing).
- Used for periodic real time tasks, eg. Video encoding/decoding.
Sched class: Real-time (RT)

- POSIX “real-time” tasks.
- Task priorities: 0 - 99.
- Inverted priorities: 0 highest in kernel, lowest in user space.
- Scheduling policies for tasks at same priority:
  - SCHED_FIFO
  - SCHED_RR, 100ms timeslice by default.
- Implemented with Linked lists.
- Used for short latency sensitive tasks, eg. IRQ threads.
Sched class: CFS (Completely fair scheduler)

- Introduced by Ingo Molnar (motivated by Rotating Staircase Deadline Scheduler by Con Kolivas).
- Scheduling policies:
  - SCHED_NORMAL: Normal Unix tasks
  - SCHED_BATCH: Batch (non-interactive) tasks
  - SCHED_IDLE: Low priority tasks
- Implemented with red-black tree (self balancing).
- Tracks Virtual runtime of tasks (amount of time task has run).
- Tasks with shortest vruntime runs first.
- Priority is used to set task’s weight, that affects vruntime.
- Higher the weight, slower will vruntime increase.
- Task’s priority is calculated by 120 + nice (-20 to +19).
- Used for all other system tasks, eg: shell.
Sched class: Idle

- Lowest priority scheduling class.
- No scheduling policies.
- One kernel thread (idle) per CPU: “swapper/N”.
- Idle thread runs only when nothing else is runnable on a CPU.
- Idle thread may take the CPU to low power state.
Runqueue

● Each CPU has an instance of “struct rq”.
● Each “rq” has DL, RT and CFS runqueues within it.
● Runnable tasks are enqueued on those runqueues.
● Lots of other information, stats are available in struct rq.

```
struct rq {
    ...
    struct cfs_rq cfs;
    struct rt_rq rt;
    struct dl_rq dl;
    ...
}
```
Runqueue (Cont.)

RQ CPU 0
- STOP
- DEADLINE
- REALTIME
- CFS
- IDLE

RQ CPU 1
- STOP
- DEADLINE
- REALTIME
- CFS
- IDLE
Thank You!!

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